

Data & Storage Services



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High orbit flight over the XRootD protocol

- Key points of the implementation
- Issues
- Status
- What's next?



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- Stateless vs. Stateful
 - Stateless: request-response (mkdir, rm, query...)
 - Stateful: open-do stuff-close (read, write, readv...)
 - The main difference in the error recovery:
 - retry or just go and ask elsewhere
 - vs. reestablish the state, possibly on another machine, and retry (not always possible)
- Solicited vs. Unsolicited
 - Responses to requests (ping-pong, classical client-server system)
 - Messages that have not been requested but affect operation (go away!, don't bother me until further notice, from now on start asking this other guy)



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- The server may respond to a request or just talk unasked:
 - **kXR_attn** unsolicited server info
 - kXR_ok the query succeeded and this is the answer
 - kXR_error failure
 - kXR_oksofar kind of okay, some part of the response arrived, but problems still may occur
 - kXR_redirect go and ask this other guy (if the request was stateful, the state needs to be re-established)
 - kXR_wait busy, bother me again in n seconds
 - kXR_waitresp the answer is not available now, will be sent in around n seconds as an unsolicited message



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- Many messages may be sent over one socket without waiting for a response and responses may arrive in any order
 - each request has a stream identifier field assigned by the client which is echoed in the response, so that the responses may be associated with the requests
- Responses to some messages may be sent over other sockets associated with the same session
 - data-heavy messages: reads and writes
 - to optimize TCP congestion avoidance/window/slowrestart issues





- stateless requests may be handled asynchronously, not only reads and writes
 - listing of huge directories an order of magnitude faster
 - and, probably, heavier for the server to handle :(
- callback model instead of request-and-poll-the-cache model
- no need to have a cache to handle async communication
- synchronous requests implemented in terms of asynchronous (with a semaphore)
- Thread safe
 - the user API classes hold very little or no mutable state at all



- one extra thread to handle socket events
- one extra thread to handle time events
- no need to spawn extra thread for every new connection
- uses host system optimized polling (through libevent) instead of block+timeout model, which should reduce number of syscalls



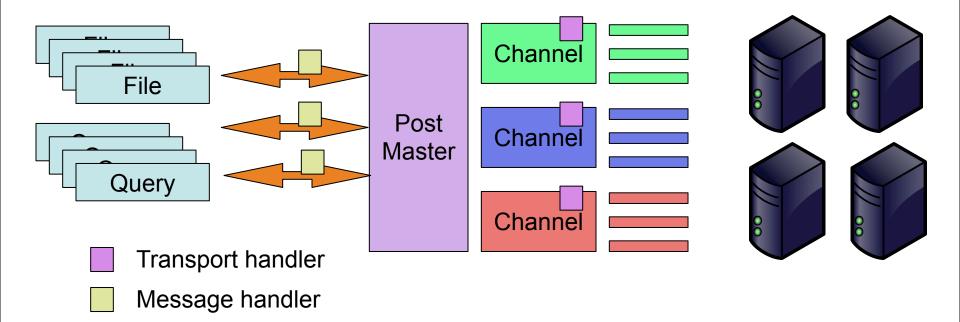


- **XrdClient::Query** for stateless requests
 - mkdir, rmdir, query, locate, move truncate, chmod, ping, stat...
- XrdClient::File for (stateful) file operations
 - read, write, readv...

User classes



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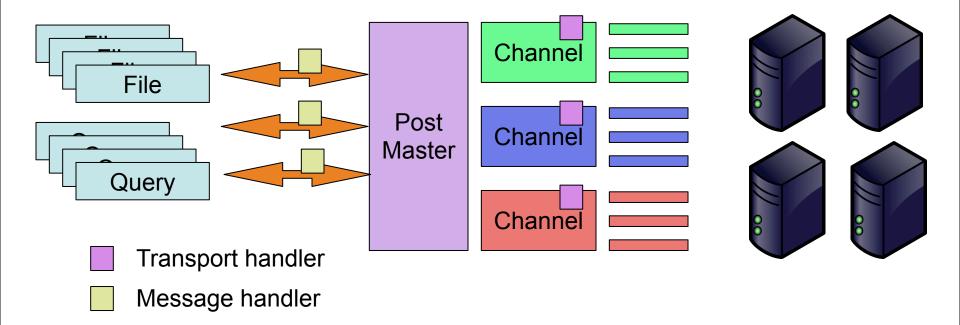


File and Query objects

- the user API entry points, they build and send the requests using a **PostMaster** object
- hold as little mutable state as possible
- register message handlers to interpret the incoming messages



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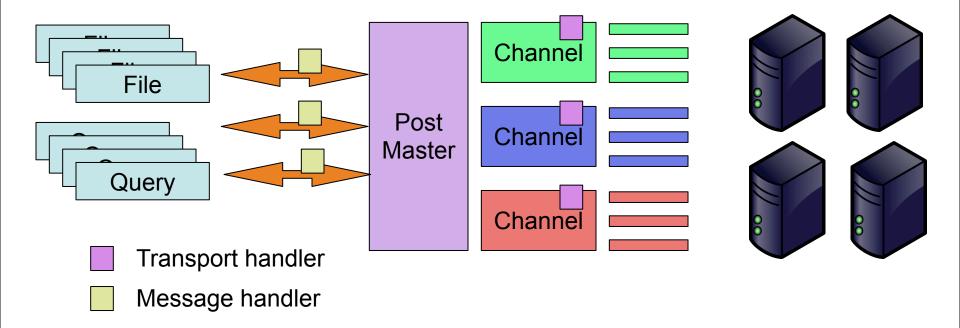


Message handlers

- receive the notification of incoming messages from the PostMaster
- handle waiting, redirection, glueing partial responses, building response objects, notifying users about final results



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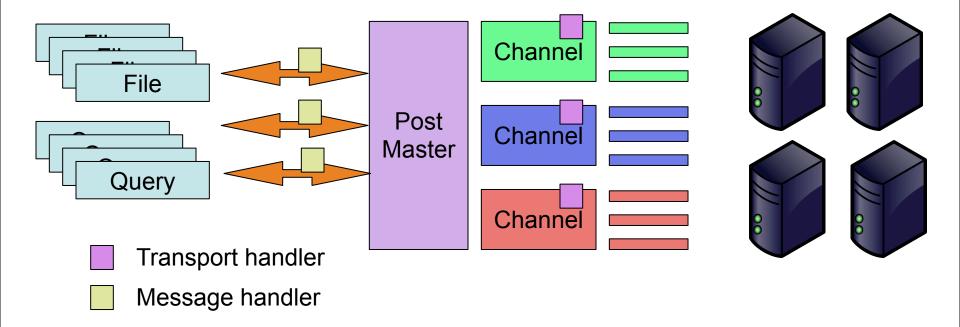


Post Master

- protocol independent (may be used for any messagebased protocol over a TCP)
- manages physical channels and streams to the servers
- sends messages to given URLs



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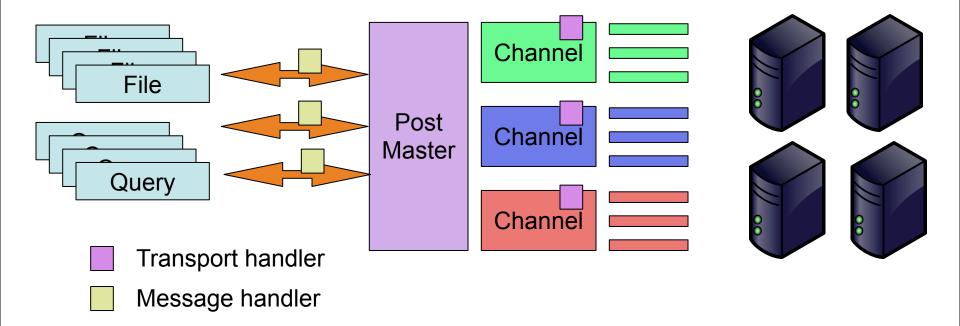


Post Master

- Manages message handler subscriptions to channels
- Notifies the handler of incoming messages



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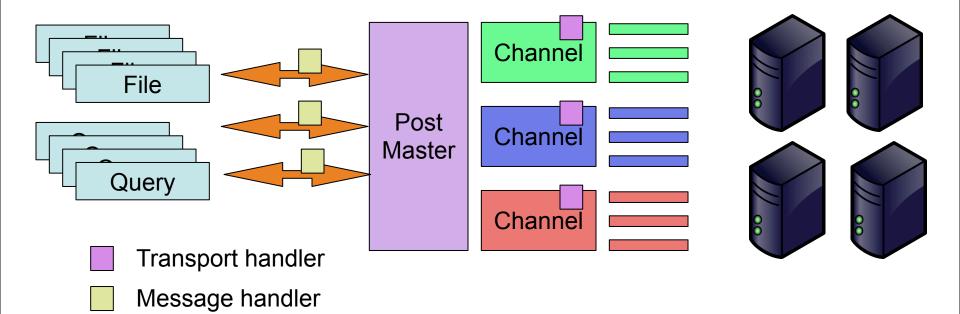


Transport handler

- entity within the post master responsible for handling specific protocol
- assigned to a channel after interpreting the protocol part of the URL (root://, xroot://, whatever)



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Transport handler

- Handles logging-in and authentications
- Stream and message multiplexing
- Extracting messges from the stream TCP socket
- Message marshalling and un-marshalling



- Stream Control Transmission Protocol
 - a transport layer protocol
 - message oriented (like UDP)
 - reliable & congestion-aware (like TCP)
 - supports multi-streaming
 - implemented in Linux and Solaris > 10
- PostMaster could be dramatically simplified if we used SCTP instead of TCP
- To be investigated

Services





- Sometimes not really clear how to deal with timeouts
 - What if the link is fine but servers tells us to wait longer than we want to?
- Recovery scenarios for complex (federated) setups.
 - Which manager to go to after encountering certain kinds of errors? Which errors should be revered in the local cluster and which on the federation level?
- How to handle write errors & redirections?
 - Start from the middle of a file on another server?





- Minor protocol issues discovered (easily fixable)
- Using libevent may be a deployment issue
 - Linux distros commonly provide 1.x but 2.x is required by the client.
- Code licensing, CERN's preference is (L)GPL, US Department of Energy prefers BSD.
- All of these to be discussed with the rest of the collaboration.





- Stateful and stateless user API for most of the requests defined by the protocol: DONE
- xrdquery utility (replacing xrd): DONE
- xrdcopy utility (replacing xrdcp): IN PROGRESS
- Unit tests: DONE
- Doxygen API documentation: DONE
- Total line count: 17274 and growing (including comments and tests)





- Caching & prefetching support
 - probably use one of existing solutions
- ROOT plugin (replacing current TXNetFile)
- IPv6 the protocol is IPv6 aware, the implementation not necessarily
- Test, test, test, test & iron-out the rough edges
 - use internally in EOS
 - planned to use in CASTOR when migrating to pure XRootD internal data transfers



Questions



Thanks for your attention!

Questions? Comments?