Scalla: Structured Clustered Architecture for Low Latency Access

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Outline



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Motivation



Initial problem: high-energy physics data in 2002

- Scalability problems with ODBMS (@1PB) High concurrent load (1000s) Large data set (2 \times disk size)
- Unnecessary DBMS overhead

Wishlist

- "Simple," file-access
 High-throughput, WAN capable
 Mostly read-only (15% create/write)
 MSS integrated
- Simplified admin → Auto fault-recovery, cluster-configuration
- Bottleneck free →Burned by lock collisions, metadata servers

General need for distributed, load-balanced tertiary-backed file access

What is Scalla



Scalla:

- a file access system built as a tree of redirectors and servers
- $\begin{tabular}{ll} \bullet & Decentralized \\ \begin{tabular}{ll} \bullet & scalabilty, no bottlenecks \\ \begin{tabular}{ll} \bullet & scalability, no bottlenecks \\ \begin{tabular}$
- Fault-tolerant → auto-mgmt handles failures and overloads
- Almost stateless

Cache-only \rightarrow cheap state changes No state for un-requested files

Minimizes centralized load (redirector)

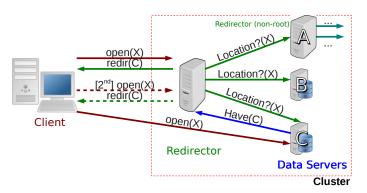
Single purpose: redirect

...up to 1000 $\frac{redirects}{sec}$ w/ 400MHz UltraSparcII ...using special data structures

Cluster architecture



64-ary trees of redirectors, servers(leaves)



N servers $\rightarrow \approx \frac{N}{64}$ redirectors and $\lceil \log_{64} N \rceil$ hops



All lookups start at redirector → redirector must be fast

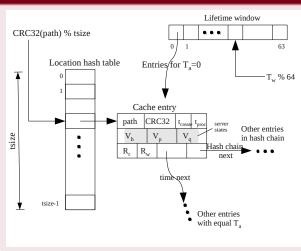
Location cache is the secret sauce

- Leverage past lookups
- Time-based eviction → bounded staleness
- Bit vectors → CPU/cache efficiency
- \bullet Lock contention/minimization, some lock-free ops \to low-latency, high concurrency
- Integration of lookup, server/location state, aging → efficiency

Caching structure



Location cache with time-based eviction



 V_h : has

 V_p : in prep

 V_q : need query

 R_r : idx read resp

 R_w : idx write resp

 T_w : clock $(\Delta t = \frac{L_t}{64})$

 t_{create} : creation (Δt)

 T_a : $t_{create}\%64$

Leverage word-length bit vectors to minimize cost, maximize cache friendliness



Idea: Tree flood, ACK-only

(efficiently query many servers)

- "request, rarely-respond"
- All servers searched in $\lceil \log_{64} n \rceil$ hops (3 hops \rightarrow 256K)
- Extremely small lookup cost per-hop
- Simplified file state: down = overloaded → not available
 Creation: add'l delay (5 seconds) req'd to avoid collision



Idea: Client-directed refresh/avoidance

(reports stale info to the redirector)

- 1 Client redirected to server
- 2 Client receives denial (or timeout) from server X
- 3 Client requests new lookup from redirector excluding X

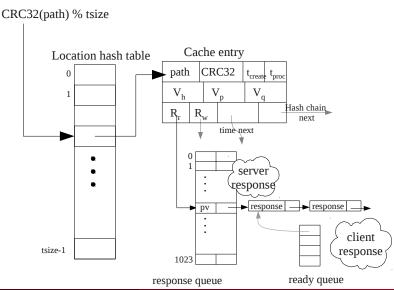
Idea: Lookup coalescing

(prevents multiple flood queries for same file)

- Using response queues $(R_r \text{ and } R_w \text{ in cache entry})$
- Processing deadlines in cache entries

Lookup response processing





Usage: File service



Widely deployed w/ original + 3rd party impl

- FGRST: @SLAC 1.6PB, 47 servers, 2 redirectors, rw)
- BaBar: 3 clusters, 75 servers, 5 redirectors, 1.5PB
- LHC

ALICE: 43 sites, ≈200 servers ATLAS production: 12 US sites

@SLAC 2.2PB, 3k-4k client jobs, 25 servers,

US federated storage: 19PB

STAR @ Brookhaven: 456 nodes

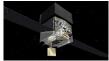










Image credits: NASA, SLAC, CERN, ATLAS Experiment (c)2012 CERN, BNL

Usage: Distributed query dispatch

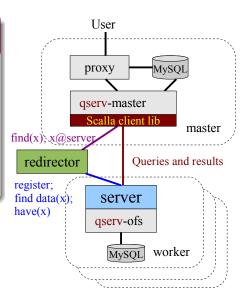


LSST's Qserv: distributed database

- data partitions → file namespace
- Queries, results →open/write/read/close
- Reuse fault-tolerant, mature distributed framework
- Tested: 150 servers, 1+2 redirectors, burst up to 9000



Image credit: LSST Corp / NOAO



Tradeoffs



- Not optimized for home directories
 - No built-in file listing, except with FUSE
- Thicker clients
 - Client mediates lookup and connections w/servers.
 - Most processing: clients + leaf servers (minimal centralization)
 - Complex client protocol (but small enough for feature-phone)
- Decentralized architecture
 - Difficult to gather system state, statistics
 - Neighbor-only connection

Summary



- Scalla provides scalable fault-tolerant file service
- Efficiency provided by integrating cluster mgmt w/ location cache
- Providing HEP/astro/physics file access for 10 years
- Providing clustering/messaging for distributed db (Qserv)

Thank you!



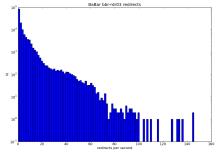
Questions?

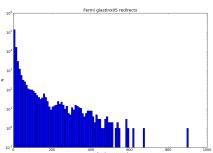
- download, refs, slides:
 - http://xrootd.slac.stanford.edu
- email:
 - Andy abh@slac.stanford.edu
 - Daniel danielw@slac.stanford.edu

Thanks!

A day of redirection







Example paths



LHC's ATLAS

 $root: //ccsrb15:1094 / / pnfs/in2p3.fr/data/atlas/atlasdatadisk/fdr08_run1/AOD/fdr08_run1.0003050....$

Qserv path examples

- Sharded query write
 xroot://qsm@mgr:1094//q/rplante_PT1_2_u_pt12prod_im3000/7505
- Sharded result readxroot://qsm@172.23.36.70:1094//result/10dbdd8da1a39908cd12b529fd79a7c4